

Multivariable knot polynomials, the V_n -polynomials, and their patterns

Shana Y. Li

University of Illinois, Urbana-Champaign

April 2025

Joint work with Stavros Garoufalidis

Table of Contents

- 1 Garoufalidis & Kashaev's multivariable knot polynomials
- 2 Computation of the V_n -polynomials
- 3 Patterns of the V_n -polynomials

Table of Contents

- 1 Garoufalidis & Kashaev's multivariable knot polynomials
- 2 Computation of the V_n -polynomials
- 3 Patterns of the V_n -polynomials

Theorem (Kashaev, 2019)

Let R be a rigid R -matrix, then the corresponding Reshetikhin–Turaev functor gives an $\text{End}(V)$ -valued invariant of oriented knots.

Theorem (Kashaev, 2019)

Let R be a rigid R -matrix, then the corresponding Reshetikhin–Turaev functor gives an $\text{End}(V)$ -valued invariant of oriented knots.

- Rigid R -matrix: an element R in $\text{Aut}(V \otimes V)$ satisfying:

Theorem (Kashaev, 2019)

Let R be a rigid R -matrix, then the corresponding Reshetikhin–Turaev functor gives an $\text{End}(V)$ -valued invariant of oriented knots.

- Rigid R -matrix: an element R in $\text{Aut}(V \otimes V)$ satisfying:
 - Yang–Baxter equation:
$$(R \otimes I) \circ (I \otimes R) \circ (R \otimes I) = (I \otimes R) \circ (R \otimes I) \circ (I \otimes R).$$

Theorem (Kashaev, 2019)

Let R be a rigid R -matrix, then the corresponding Reshetikhin–Turaev functor gives an $\text{End}(V)$ -valued invariant of oriented knots.

- Rigid R -matrix: an element R in $\text{Aut}(V \otimes V)$ satisfying:
 - Yang–Baxter equation:
$$(R \otimes I) \circ (I \otimes R) \circ (R \otimes I) = (I \otimes R) \circ (R \otimes I) \circ (I \otimes R).$$
 - Rigidity: the *partial transposes*
$$\widetilde{R^{\pm 1}} := (\varepsilon \otimes I \otimes I) \circ (I \otimes R^{\pm 1} \otimes I) \circ (I \otimes I \otimes \eta)$$
 are invertible.
 $\varepsilon: V \otimes V \rightarrow \mathbb{F}$ and $\eta: \mathbb{F} \rightarrow V \otimes V$: the evaluation and coevaluation maps.

Theorem (Kashaev, 2019)

Let R be a rigid R -matrix, then the corresponding Reshetikhin–Turaev functor gives an $\text{End}(V)$ -valued invariant of oriented knots.

- Rigid R -matrix: an element R in $\text{Aut}(V \otimes V)$ satisfying:
 - Yang–Baxter equation:
$$(R \otimes I) \circ (I \otimes R) \circ (R \otimes I) = (I \otimes R) \circ (R \otimes I) \circ (I \otimes R).$$
 - Rigidity: the *partial transposes*
$$\widetilde{R^{\pm 1}} := (\varepsilon \otimes I \otimes I) \circ (I \otimes R^{\pm 1} \otimes I) \circ (I \otimes I \otimes \eta)$$
 are invertible.
 $\varepsilon: V \otimes V \rightarrow \mathbb{F}$ and $\eta: \mathbb{F} \rightarrow V \otimes V$: the evaluation and coevaluation maps.
- Reshetikhin–Turaev functor: a functor (determined by R) from the category of tangles to the category of vector spaces.

Theorem (Garoufalidis & Kashaev, 2023)

Given a braided Hopf algebra with automorphisms, one can construct a rigid R -matrix.

Theorem (Garoufalidis & Kashaev, 2023)

Given a braided Hopf algebra with automorphisms, one can construct a rigid R -matrix.

The procedure:

$$\left\{ \begin{array}{l} \text{Braided} \\ \text{Hopf algebras} \\ \text{with autos} \end{array} \right\} \rightarrow \left\{ \begin{array}{l} \text{Braided} \\ \text{Yetter-Drinfel'd} \\ \text{modules with autos} \end{array} \right\} \rightarrow \left\{ \begin{array}{l} \text{Rigid} \\ R\text{-matrices} \end{array} \right\}$$

Summary: a systematic source of knot invariants:

$$\left\{ \begin{array}{l} \text{Braided} \\ \text{Hopf algebras} \\ \text{with autos} \end{array} \right\} \xrightarrow{\text{G \& K, 2023}} \left\{ \begin{array}{l} \text{Rigid} \\ R\text{-matrices} \end{array} \right\} \xrightarrow{\text{K, 2019}} \{\text{Knot invariants}\}$$

Summary: a systematic source of knot invariants:

$$\left\{ \begin{array}{l} \text{Braided} \\ \text{Hopf algebras} \\ \text{with autos} \end{array} \right\} \xrightarrow{\text{G \& K, 2023}} \left\{ \begin{array}{l} \text{Rigid} \\ R\text{-matrices} \end{array} \right\} \xrightarrow{\text{K, 2019}} \{\text{Knot invariants}\}$$

One source of braided Hopf algebras: Nichols algebras.

Summary: a systematic source of knot invariants:

$$\left\{ \begin{array}{l} \text{Braided} \\ \text{Hopf algebras} \\ \text{with autos} \end{array} \right\} \xrightarrow{\text{G \& K, 2023}} \left\{ \begin{array}{l} \text{Rigid} \\ R\text{-matrices} \end{array} \right\} \xrightarrow{\text{K, 2019}} \{\text{Knot invariants}\}$$

One source of braided Hopf algebras: Nichols algebras.

- Nichols algebras of rank 1: recovers the colored Jones polynomials and the ADO polynomials.

Summary: a systematic source of knot invariants:

$$\left\{ \begin{array}{l} \text{Braided} \\ \text{Hopf algebras} \\ \text{with autos} \end{array} \right\} \xrightarrow{\text{G \& K, 2023}} \left\{ \begin{array}{l} \text{Rigid} \\ R\text{-matrices} \end{array} \right\} \xrightarrow{\text{K, 2019}} \{\text{Knot invariants}\}$$

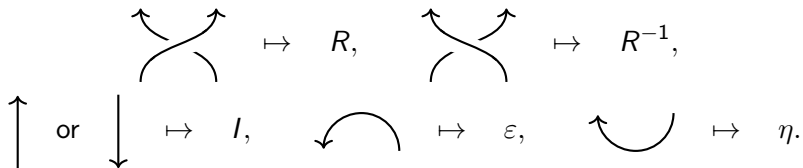
One source of braided Hopf algebras: Nichols algebras.

- Nichols algebras of rank 1: recovers the colored Jones polynomials and the ADO polynomials.
- Nichols algebras of rank 2: recovers the Links–Gould polynomial, and gives the V_n -polynomials.

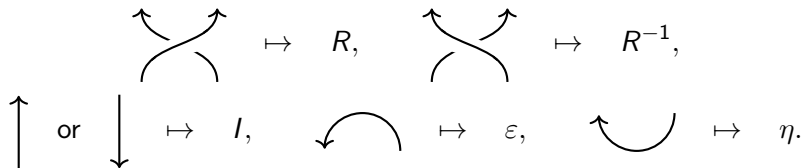
Table of Contents

- 1 Garoufalidis & Kashaev's multivariable knot polynomials
- 2 Computation of the V_n -polynomials
- 3 Patterns of the V_n -polynomials

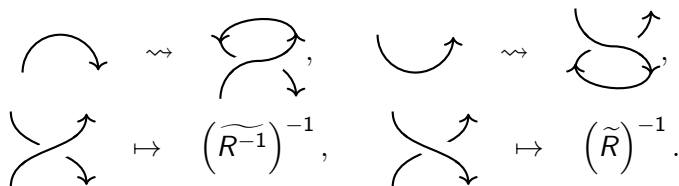
Reshetikhin–Turaev functor:



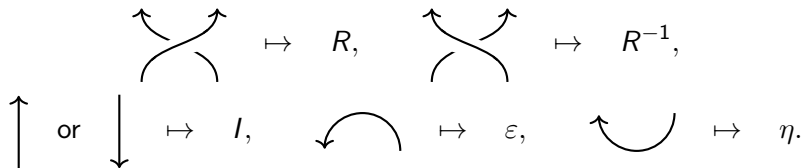
Reshetikhin–Turaev functor:



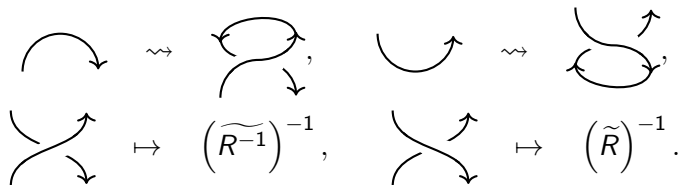
For local extrema going from left to right:



Reshetikhin–Turaev functor:

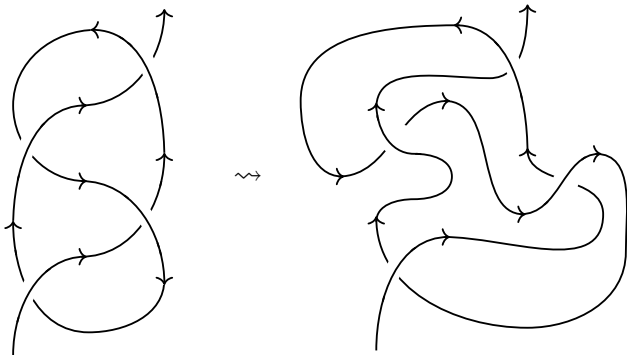


For local extrema going from left to right:

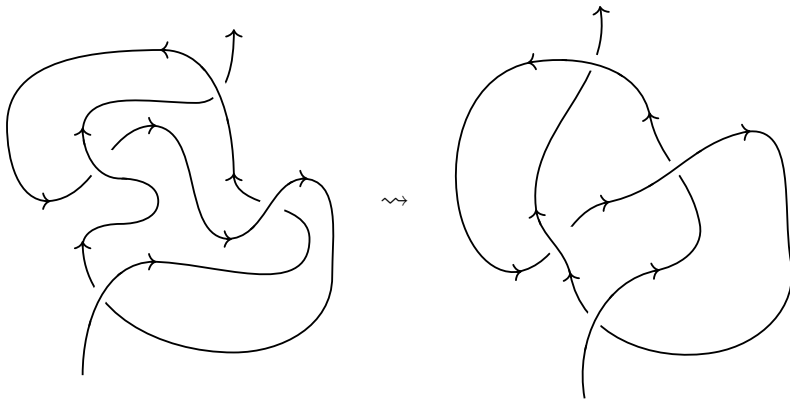


For V_n -polynomials, $\varepsilon \circ (\widetilde{R^{-1}})^{-1} = (\widetilde{R^{-1}})^{-1} \circ \eta = \varepsilon \circ (\widetilde{R})^{-1} = (\widetilde{R})^{-1} \circ \eta$ is a diagonalizable matrix with only ± 1 's on the diagonal.

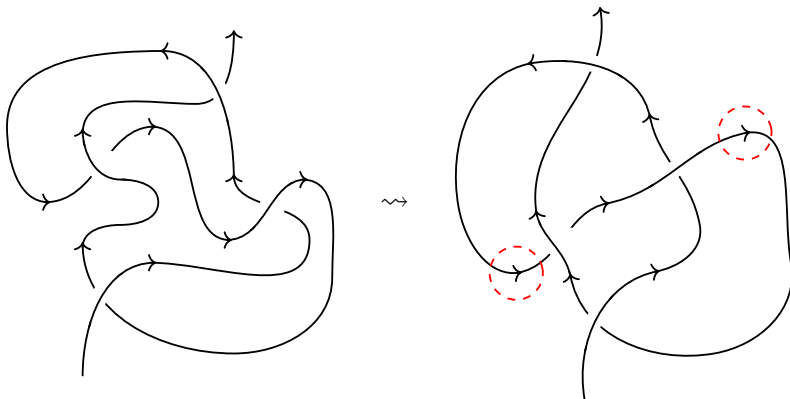
Example: the 4_1 knot



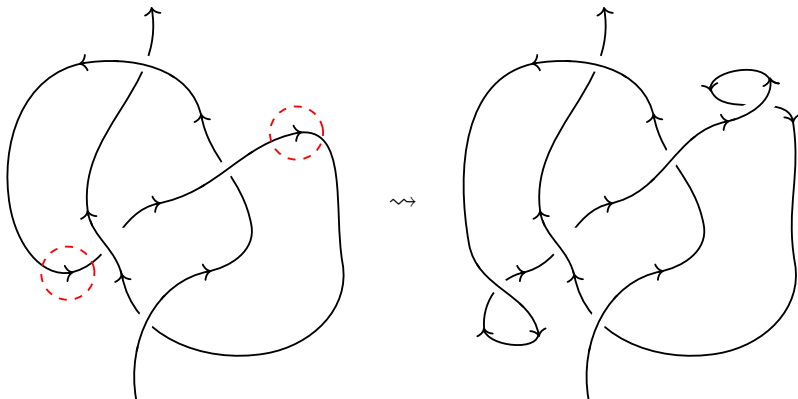
Example: the 4_1 knot



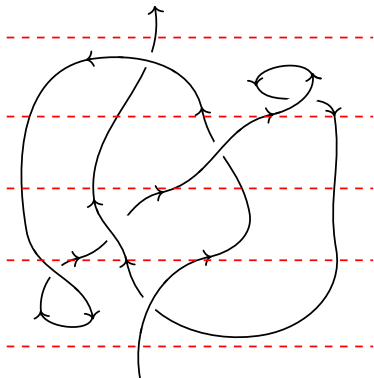
Example: the 4_1 knot



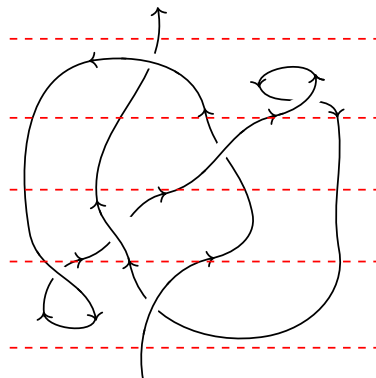
Example: the 4_1 knot



Example: the 4_1 knot



$$\begin{array}{r}
 V \\
 \uparrow \\
 V^{\otimes 5} \\
 \uparrow \\
 V^{\otimes 5} \\
 \uparrow \\
 V^{\otimes 5} \\
 \uparrow \\
 V
 \end{array}
 \begin{array}{l}
 \\
 \left((\varepsilon \otimes I) \circ (I \otimes R^{-1}) \right) \otimes \left(\varepsilon \circ (\widetilde{R})^{-1} \right) \\
 \\
 I \otimes I \otimes R \otimes I \\
 \\
 I \otimes R^{-1} \otimes I \otimes I \\
 \\
 \left((\widetilde{R^{-1}})^{-1} \circ \eta \right) \otimes ((R \otimes I) \circ (I \otimes \eta)) \\
 \\
 \end{array}$$

Example: the 4_1 knot

$$\begin{array}{c}
 V \\
 \uparrow \\
 V^{\otimes 5} \\
 \uparrow \\
 V^{\otimes 5} \\
 \uparrow \\
 V^{\otimes 5} \\
 \uparrow \\
 V
 \end{array}
 \begin{array}{l}
 \\
 \left((\varepsilon \otimes I) \circ (I \otimes R^{-1}) \right) \otimes \left(\varepsilon \circ (\tilde{R})^{-1} \right) \\
 \\
 I \otimes I \otimes R \otimes I \\
 \\
 I \otimes R^{-1} \otimes I \otimes I \\
 \\
 \left((\widetilde{R^{-1}})^{-1} \circ \eta \right) \otimes \left((R \otimes I) \circ (I \otimes \eta) \right) \\
 \\
 \\
 \end{array}
 \begin{array}{l}
 \\
 \\
 \\
 \mapsto \\
 \\
 \\
 \\
 \\
 \\
 \\
 \end{array}$$

Fact

V is a submodule of an algebra over \mathbb{F} , and $1 \in V$ is an eigenvector of the above $\text{End}(V)$ -valued invariant.

Fix a basis $\mathcal{B} := \{e_i\}$ of V , $R^{\pm 1}$ become matrices whose entries can be denoted by $(R^{\pm 1})_{e_i \otimes e_j}^{e_k \otimes e_l}$.

Fix a basis $\mathcal{B} := \{e_i\}$ of V , $R^{\pm 1}$ become matrices whose entries can be denoted by $(R^{\pm 1})_{e_i \otimes e_j}^{e_k \otimes e_l}$.

To compute the eigenvalue of the $\text{End}(V)$ -valued invariant is to evaluate a sum of the form

$$\sum_{\substack{a_1, \dots, a_{2c-1} \in \mathcal{B} \\ a_0 = a_{2c} = 1}} \pm \underbrace{(R^{\pm 1})_{a_0 \otimes a_1}^{a_2 \otimes a_3} \cdots (R^{\pm 1})_{a_{2c-3} \otimes a_{2c-2}}^{a_{2c-1} \otimes a_{2c}}}_{\text{a product of length } c},$$

where c is the number of crossings of the knot. This sum is called the *state sum*.

Fix a basis $\mathcal{B} := \{e_i\}$ of V , $R^{\pm 1}$ become matrices whose entries can be denoted by $(R^{\pm 1})_{e_i \otimes e_j}^{e_k \otimes e_l}$.

To compute the eigenvalue of the $\text{End}(V)$ -valued invariant is to evaluate a sum of the form

$$\sum_{\substack{a_1, \dots, a_{2c-1} \in \mathcal{B} \\ a_0 = a_{2c} = 1}} \pm \underbrace{(R^{\pm 1})_{a_0 \otimes a_1}^{a_2 \otimes a_3} \cdots (R^{\pm 1})_{a_{2c-3} \otimes a_{2c-2}}^{a_{2c-1} \otimes a_{2c}}}_{\text{a product of length } c},$$

where c is the number of crossings of the knot. This sum is called the *state sum*.

Therefore, it requires

$$c \cdot (\dim V)^{2c-1}$$

times of computations to compute the eigenvalue.

For V_n -polynomials, $\dim V = 4n$.

For V_n -polynomials, $\dim V = 4n$.

With $n = 2$, for the simplest knot 3_1 , we have

$$c \cdot (\dim V)^{2c-1} = 98,304.$$

For V_n -polynomials, $\dim V = 4n$.

With $n = 2$, for the simplest knot 3_1 , we have

$$c \cdot (\dim V)^{2c-1} = 98,304.$$

This is how you know that you wouldn't want to compute it by hand.

For V_n -polynomials, $\dim V = 4n$.

With $n = 2$, for the simplest knot 3_1 , we have

$$c \cdot (\dim V)^{2c-1} = 98,304.$$

This is how you know that you wouldn't want to compute it by hand.

For $c = 12$,

For V_n -polynomials, $\dim V = 4n$.

With $n = 2$, for the simplest knot 3_1 , we have

$$c \cdot (\dim V)^{2c-1} = 98,304.$$

This is how you know that you wouldn't want to compute it by hand.

For $c = 12$,

$$c \cdot (\dim V)^{2c-1} = 7,083,549,724,304,467,820,544,$$

For V_n -polynomials, $\dim V = 4n$.

With $n = 2$, for the simplest knot 3_1 , we have

$$c \cdot (\dim V)^{2c-1} = 98,304.$$

This is how you know that you wouldn't want to compute it by hand.

For $c = 12$,

$$c \cdot (\dim V)^{2c-1} = 7,083,549,724,304,467,820,544,$$

and for $c = 15$,

For V_n -polynomials, $\dim V = 4n$.

With $n = 2$, for the simplest knot 3_1 , we have

$$c \cdot (\dim V)^{2c-1} = 98,304.$$

This is how you know that you wouldn't want to compute it by hand.

For $c = 12$,

$$c \cdot (\dim V)^{2c-1} = 7,083,549,724,304,467,820,544,$$

and for $c = 15$,

$$c \cdot (\dim V)^{2c-1} = 2,321,137,573,660,088,015,435,857,920.$$

For V_n -polynomials, $\dim V = 4n$.

With $n = 2$, for the simplest knot 3_1 , we have

$$c \cdot (\dim V)^{2c-1} = 98,304.$$

This is how you know that you wouldn't want to compute it by hand.

For $c = 12$,

$$c \cdot (\dim V)^{2c-1} = 7,083,549,724,304,467,820,544,$$

and for $c = 15$,

$$c \cdot (\dim V)^{2c-1} = 2,321,137,573,660,088,015,435,857,920.$$

Worse, the entries $(R^{\pm 1})_{e_i \otimes e_j}^{e_k \otimes e_l}$ are polynomials in two variables, instead of scalars.

For V_n -polynomials, $\dim V = 4n$.

With $n = 2$, for the simplest knot 3_1 , we have

$$c \cdot (\dim V)^{2c-1} = 98,304.$$

This is how you know that you wouldn't want to compute it by hand.

For $c = 12$,

$$c \cdot (\dim V)^{2c-1} = 7,083,549,724,304,467,820,544,$$

and for $c = 15$,

$$c \cdot (\dim V)^{2c-1} = 2,321,137,573,660,088,015,435,857,920.$$

Worse, the entries $(R^{\pm 1})_{e_i \otimes e_j}^{e_k \otimes e_l}$ are polynomials in two variables, instead of scalars. We computed the V_2 -polynomials for all knots with ≤ 15 crossings, and more.

To optimize the computation:

To optimize the computation:

- The R -matrix is sparse: a divide and conquer method sees the 0's at each step and eliminates a lot of terms.

n	Nonzero elements (%)	# R
2	177 (4.3%)	4096
3	585 (2.8%)	20,736
4	1377 (2.1%)	65,536

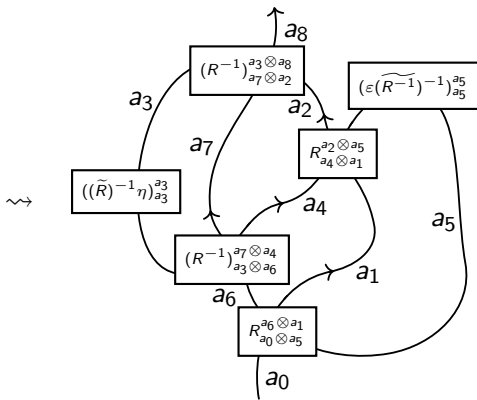
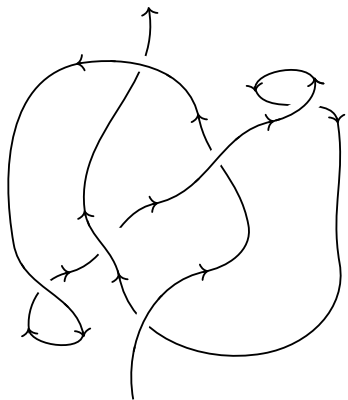
To optimize the computation:

- The R -matrix is sparse: a divide and conquer method sees the 0's at each step and eliminates a lot of terms.

n	Nonzero elements (%)	# R
2	177 (4.3%)	4096
3	585 (2.8%)	20,736
4	1377 (2.1%)	65,536

- Use optimized tensor contraction path.

Example: the 4_1 knot again



Example: the 4_1 knot again

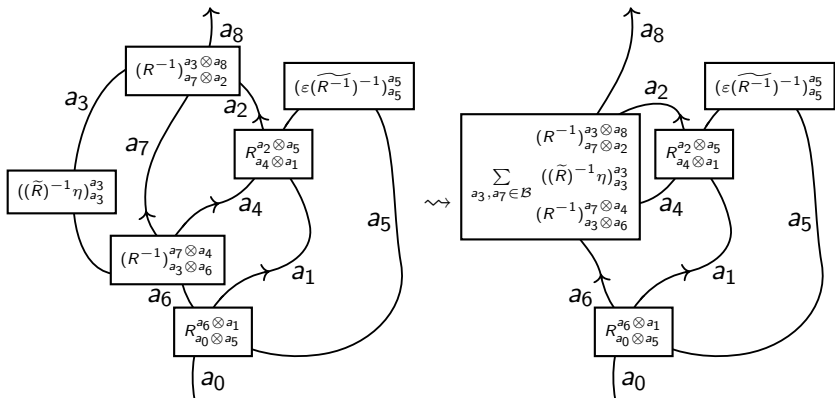


Table of Contents

- 1 Garoufalidis & Kashaev's multivariable knot polynomials
- 2 Computation of the V_n -polynomials
- 3 Patterns of the V_n -polynomials

Let $V_{K,n}(t, q) \in \mathbb{Z}[t^{\pm 1}, q^{\pm 1}]$ be the V_n -polynomial of knot K in variables t and q .

Let $V_{K,n}(t, q) \in \mathbb{Z}[t^{\pm 1}, q^{\pm 1}]$ be the V_n -polynomial of knot K in variables t and q .

- Symmetry:

$$V_{K,n}(t, q) = V_{K,n}(t^{-1}, q), \quad V_{\overline{K},n}(t, q) = V_{K,n}(t, q^{-1})$$

Let $V_{K,n}(t, q) \in \mathbb{Z}[t^{\pm 1}, q^{\pm 1}]$ be the V_n -polynomial of knot K in variables t and q .

- Symmetry:

$$V_{K,n}(t, q) = V_{K,n}(t^{-1}, q), \quad V_{\overline{K},n}(t, q) = V_{K,n}(t, q^{-1})$$

- Specialization (conjecturally):

$$V_{K,n}(q^{n/2}, q) = 1, \quad V_{K,n}(t, 1) = \Delta_K(t)^2$$

where $\Delta_K(t)$ is the Alexander polynomial.

Let $V_{K,n}(t, q) \in \mathbb{Z}[t^{\pm 1}, q^{\pm 1}]$ be the V_n -polynomial of knot K in variables t and q .

- Symmetry:

$$V_{K,n}(t, q) = V_{K,n}(t^{-1}, q), \quad V_{\overline{K},n}(t, q) = V_{K,n}(t, q^{-1})$$

- Specialization (conjecturally):

$$V_{K,n}(q^{n/2}, q) = 1, \quad V_{K,n}(t, 1) = \Delta_K(t)^2$$

where $\Delta_K(t)$ is the Alexander polynomial.

- Genus bound (conjecturally):

$$\deg_t V_{K,n}(t, q) \leq 4g(K)$$

where $g(K)$ is Seifert genus of K .

Theorems

- (GKKST) The V_1 -polynomial is the Links–Gould polynomial.
- (KT) The Links–Gould polynomial satisfies both the specialization and the genus bound.
- $V_{K,2}$ is determined by

$$V_{K,2}(t^2, q^2) = c_{2,0}(t, q)V_{K(2,1),1}(t, q) + c_{2,-1}(t, q)V_{K,1}(t^2q^{-1}, q) + c_{2,1}(t, q)V_{K,1}(t^2q, q)$$

where $K(2, 1)$ is the $(2, 1)$ -parallel of K .

Theorems

- (GKKST) The V_1 -polynomial is the Links–Gould polynomial.
- (KT) The Links–Gould polynomial satisfies both the specialization and the genus bound.
- $V_{K,2}$ is determined by

$$V_{K,2}(t^2, q^2) = c_{2,0}(t, q)V_{K(2,1),1}(t, q) + c_{2,-1}(t, q)V_{K,1}(t^2q^{-1}, q) + c_{2,1}(t, q)V_{K,1}(t^2q, q)$$

where $K(2, 1)$ is the $(2, 1)$ -parallel of K .

Since $g(K(2, 1)) = 2g(K)$, the last statement implies that V_2 also satisfies both the specialization and the genus bound.

Theorems

- (GKKST) The V_1 -polynomial is the Links–Gould polynomial.
- (KT) The Links–Gould polynomial satisfies both the specialization and the genus bound.
- $V_{K,2}$ is determined by

$$V_{K,2}(t^2, q^2) = c_{2,0}(t, q)V_{K(2,1),1}(t, q) + c_{2,-1}(t, q)V_{K,1}(t^2q^{-1}, q) + c_{2,1}(t, q)V_{K,1}(t^2q, q)$$

where $K(2, 1)$ is the $(2, 1)$ -parallel of K .

Since $g(K(2, 1)) = 2g(K)$, the last statement implies that V_2 also satisfies both the specialization and the genus bound.

Conjecturally, V_n -polynomials satisfy relations similar to the one above.

Question

When is the equality achieved in the genus bound inequality?

$$\deg_t V_{K,n}(t, q) \leq 4g(K)$$

Question

When is the equality achieved in the genus bound inequality?

$$\deg_t V_{K,n}(t, q) \leq 4g(K)$$

With the specialization, we have

$$2 \deg_t \Delta_K(t) \leq \deg_t V_{K,n}(t, q) \leq 4g(K).$$

Question

When is the equality achieved in the genus bound inequality?

$$\deg_t V_{K,n}(t, q) \leq 4g(K)$$

With the specialization, we have

$$2 \deg_t \Delta_K(t) \leq \deg_t V_{K,n}(t, q) \leq 4g(K).$$

Since Alexander polynomials satisfy $\deg_t \Delta_K(t) \leq 2g(K)$, a sufficient condition:

$$\deg_t \Delta_K(t) = 2g(K). \quad (1)$$

Question

When is the equality achieved in the genus bound inequality?

$$\deg_t V_{K,n}(t, q) \leq 4g(K)$$

With the specialization, we have

$$2 \deg_t \Delta_K(t) \leq \deg_t V_{K,n}(t, q) \leq 4g(K).$$

Since Alexander polynomials satisfy $\deg_t \Delta_K(t) \leq 2g(K)$, a sufficient condition:

$$\deg_t \Delta_K(t) = 2g(K). \quad (1)$$

We call knots satisfying eq. (1) *tight*, and others *loose*.

There are no loose knots with ≤ 10 crossings.

There are no loose knots with ≤ 10 crossings.

crossings	11	12	13	14	15	16
Knots	552	2176	9988	46972	253293	1388705
Loose knots	7	29	208	1220	6319	48174

Table: Knot counts, up to mirror image

There are no loose knots with ≤ 10 crossings.

crossings	11	12	13	14	15	16
Knots	552	2176	9988	46972	253293	1388705
Loose knots	7	29	208	1220	6319	48174

Table: Knot counts, up to mirror image

polynomial	V_1	V_2	V_3	V_4
Knots	≤ 15	≤ 15	≤ 11	≤ 10
Loose knots	≤ 16	≤ 16		

Table: Computed knots for each V_n

There are no loose knots with ≤ 10 crossings.

crossings	11	12	13	14	15	16
Knots	552	2176	9988	46972	253293	1388705
Loose knots	7	29	208	1220	6319	48174

Table: Knot counts, up to mirror image

polynomial	V_1	V_2	V_3	V_4
Knots	≤ 15	≤ 15	≤ 11	≤ 10
Loose knots	≤ 16	≤ 16		

Table: Computed knots for each V_n

crossings	11	12	13	14	15	16
V_1 genus bound $<$	7	20	173	974	5025	37205
V_2 genus bound $<$	0	0	0	0	0	0

Table: Non-sharp genus bound counts

Theorem (Garoufalidis & Li, 2024)

The genus bound inequality is an equality for V_2 -polynomials for all 1,701,936 knots with ≤ 16 crossings.

Theorem (Garoufalidis & Li, 2024)

The genus bound inequality is an equality for V_2 -polynomials for all 1,701,936 knots with ≤ 16 crossings.

In other words, the V_2 -polynomials (conjecturally) detect the genus – an achievement only the knot Floer homology has made before.

Theorem (Garoufalidis & Li, 2024)

The genus bound inequality is an equality for V_2 -polynomials for all 1,701,936 knots with ≤ 16 crossings.

In other words, the V_2 -polynomials (conjecturally) detect the genus – an achievement only the knot Floer homology has made before.

Question

Does the V_2 -polynomials *actually* detect the genus of knots? Why?

Question

When do two knots have equal V_2 polynomial?

Question

When do two knots have equal V_2 polynomial?

crossings	≤ 11	12	13	14	15
pairs	0	3	25	187	2324
triples	0	0	0	1	38

Table: Number of V_2 -equivalence classes of size more than 1 (up to mirror image).

Question

When do two knots have equal V_2 polynomial?

crossings	≤ 11	12	13	14	15
pairs	0	3	25	187	2324
triples	0	0	0	1	38

Table: Number of V_2 -equivalence classes of size more than 1 (up to mirror image).

Theorem (Garoufalidis & Li, 2024)

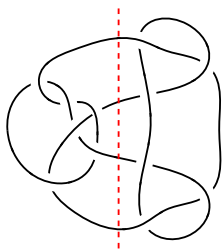
All knots with ≤ 15 crossings in the same V_2 -equivalence classes

- have equal HFK and equal Khovanov Homology,
- are Conway mutant knots to each other, in particular they have equal volumes, trace field, colored Jones polynomials, ADO polynomials and HOMFLY polynomial.

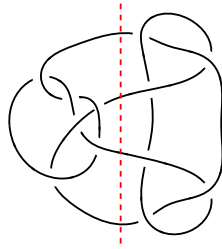
Theorem (Garoufalidis & Li, 2024)

All knots with ≤ 15 crossings in the same V_2 -equivalence classes

- have equal HFK and equal Khovanov Homology,
- are Conway mutant knots to each other, in particular they have equal volumes, trace field, colored Jones polynomials, ADO polynomials and HOMFLY polynomial.



12n364



12n365

Theorem (Garoufalidis & Li, 2024)

All knots with ≤ 15 crossings in the same V_2 -equivalence classes

- have equal HFK and equal Khovanov Homology,
- are Conway mutant knots to each other, in particular they have equal volumes, trace field, colored Jones polynomials, ADO polynomials and HOMFLY polynomial.

crossings	11	12	13	14	15
V_2 -equiv classes	0	3	25	188	2362
mutant classes	16	75	774	4435	29049

Table: Number of nontrivial V_2 -equiv classes versus Conway mutant classes.

Theorem (Garoufalidis & Li, 2024)

All knots with ≤ 15 crossings in the same V_2 -equivalence classes

- have equal HFK and equal Khovanov Homology,
- are Conway mutant knots to each other, in particular they have equal volumes, trace field, colored Jones polynomials, ADO polynomials and HOMFLY polynomial.

Question

Are V_2 -equivalent knots always Conway mutant? Do they always have equal HFK and equal Khovanov Homology? Why?

Question

Are V_2 -equivalent knots always Conway mutant? Do they always have equal HFK and equal Khovanov Homology? Why?

A partial answer: most of them are HFK-thin and Khovanov-thin, for which equal HFK and equal Khovanov Homology follows given the mutant condition.

Question

Are V_2 -equivalent knots always Conway mutant? Do they always have equal HFK and equal Khovanov Homology? Why?

A partial answer: most of them are HFK-thin and Khovanov-thin, for which equal HFK and equal Khovanov Homology follows given the mutant condition.

total	tight & thin	tight & thick	loose & thick
2578	1877	457	244

Table: Number of nontrivial V_2 -equiv classes in each flavor, up to 15 crossings.

A Conspiracy Theory:

A Conspiracy Theory:

Proposition

For all alternating knots with ≤ 15 crossings, we have

$$V_1(t, -q), V_2(t, -q) \in \mathbb{Z}_{\geq 0}[t^{\pm 1}, q^{\pm 1}].$$

A Conspiracy Theory:

Proposition

For all alternating knots with ≤ 15 crossings, we have

$$V_1(t, -q), V_2(t, -q) \in \mathbb{Z}_{\geq 0}[t^{\pm 1}, q^{\pm 1}].$$

Question

Does this indicate a categorification of V_1 and V_2 ?